RECALL OUR CONVERGENCE TESTS WITH THE QUARTER-CIRCLE INTEGRATION. WE COULD CONSTRUCT AN INTEGRATION OF A SPHERE OVER ONE OCTANT BY NESTING THESE.

$$\frac{1}{2} \qquad f(x,y) = \frac{1}{2} = \sqrt{1 - (x^2 + y^2)}$$

$$x = 0, \quad Lower \ Limit$$

$$x = \sqrt{1 - y^2}, \quad Vepger \ Limit$$

$$x = \sqrt{1 - (x^2 + y^2)}$$

$$y = 0, \quad x = 0$$

BASED ON OVA PRIOR EXPERIENCE, CONVERGENCE OF THE QUARTER-CIRCLE INTEGRAL 15

O(N-3/2) - REGARDLESS OF THE NOMINAL "ORDER" OF THE QUADRATURE.

SO OUR DOUBLE INTEGRAL WOULD CONVERGE LIKE

$$O(n^{-3/4}).$$

E.G. ACLURACY OF 10-3 REQUIRES 22 = 104.

EXAMPLE, COST.

SUPPOSE WE ARE INTEGRATING THE V-DIMENSIONAL HYPERSPHERE TO FIND ITS VOLUME.

M = V - 1

Convergence:

$$O(n^{-3/2m}) = O(n^{\frac{-3}{2(\nu-1)}})$$

NOTE HOW BAD THINGS HAVE BECOME WITH A MODERATELY "MASTY" INTEGRAND!

E.G.
$$V = 4 - D$$
 HYPERSPHERE $\rightarrow O(n^{-1/2})$
 $V = 7 - D$ HYPERSPHERE $\rightarrow O(n^{-1/4})$

Q: IF I THROW 100 DARTS AT A V-DIMENSIONAL DART BOARD AND 25 HIT THE BULLSEY, WHAT IS THE PROBABILITY THAT MY NEXT THROW WILL HIT, AND WHAT IS THE UNCERTAINTY?

$$P = \begin{cases} d^{2} \times P(\bar{x}) = ? \pm ? \\ d^{2} \times P(\bar{x}) = ? \pm ? \end{cases}$$

A: $p \approx \frac{25 \pm 5}{100} = 0.25 \pm 0.05$

a: HOW WOULD IP CONVERGE AS I INCREASE / FROM 100 TO 1,000,000? HOW DOES THIS DEPEND ON V?

| | THE DART BOARD IDEA IS CALLED MONTE CARLO INTEGRATION, (8) |
|--|--|
| , | LET'S BE MORE PRECISE. |
| | LOOK AT hypersphere 1. m & Discuss. |
| | THE ALGORITHM: $(\vec{x}', \vec{x}', \vec{x}'')$ (D) CHOOSE M. POINTS, WITHIN THE HYPERCUBE, WHERE |
| | $0 < x_n^2 < 1$, $n = 1, 2, V$, FROM UNIFORM DISTRIBUTION. |
| | 2) EVALUATE HOW MANY POINTS ARE WITHIN THE HYPGRSPHERE, M. |
| | $\langle n' \rangle = n_p$ |
| - | WHERE 1P = CHAPCE OF HITTING THE BULLSEYE (HTPERSPHERE) |
| To Sentitivities and an expense | P = VOLUME OF HYPERSPHERE VOLUME OF HYPERCUBE |
| ACADOMIC COMMISSION OF THE PARTY OF THE PART | Volume of Hypercube |
| The state of the s | (a) YOLUME OF HYPERSPHERE = $p \cdot 2^{\nu} \approx \frac{n'}{n} 2^{\nu}$ |
| The same and the s | WE EXPECT CONVERGENCE OF O (92-1/2). |

MORE CAREFUL CALCULATION OF MONTE CARLO UNCERTAINTY... THIS comes bown to Estimating p:

THE DISTRIBUTION OF M, P(M/N,P)

15 A BINOMIAL DISTAIGUTION.

A SINGLE TRIAL IS EITHER O OR 1 COUNT, WITH A MEAN OF P AND VARIANCE $\sigma_{i}^{2} = \sum_{x} p(x) (x-\bar{x})^{2} = 10 (1-p)^{2} + (1-p)(0-p)^{2} = p(1-p)$

TO OBTAIN THE VARIANCE FOR 92 TRIAZS, WE ADD IN QUADRATURE:

$$\sigma_n^2 = \mathcal{N}\sigma_1^2 = \mathcal{N}_p(1-p).$$

THUS, WHEN WE MEASURE 72, WE GET

$$n' \approx np \pm \sqrt{np(1-p)}$$

$$\Rightarrow \left(p \approx \frac{n'}{n} \pm \sqrt{\frac{p(1-p)}{n}} \right)$$

FRACTIONAL UNCERTAINTY IN P, ALO THUS

IN THE VOLUME OF THE HYPERSPAGRE, IS

$$\frac{\Delta V}{V} = \frac{\Delta P}{P} = \sqrt{\frac{(1-P)}{NP}} + \frac{AS EXPECTED}{O(n^{-1/2})}$$

THIS IS THE GENERAL FORM OF THE UNCERTAINTY FOR ANY SIMPLE MONTE CARLO CALCULATION (NOT JUST THE HYPERSPHERE VOLUME — THOSE DETAILS DID NOT ENTER IN TO OUR CALCULATION).

RANDOM NUMBER GENERATION — [NR CH.7]

RNG IS ESSENTIAL TO MONTE-CARW. OCTAVE HAS
A PRETTY GOOD RNG BUILT INTO THE RAND FUNCTION.
YOU CAN USE IT BLINDLY, BUT IT'S WORTH KNOWING
WHAT IS GOING ON UNDER THE HOOD.
IMPORTANT CONCEPTS:

- ① COMPUTERS ARE DETERMINISTIC MACHINES.

 RNG ALGORITHMS RELY ON "PSEUDO-RANDON"

 SEQUENCES" SERVENCES OF BIG INTEGERS* THAT

 APPEAR TO BE STATISTICALLY MOEPENOENT.
 - 1 ALL PSEUDO-RANDOM SEQUENCES ARE ACTUALY PERIODIC.

FOR EXAMPLE, OCTAVE'S RAND REPEATS AFTER 2 19937 - 1 CYCLES. (THAT'S ~ 10 6000 --!)

THE "STATE" OF THE RNG IS THE BIG INTEGER,
WHICH WILL BE PERMUTED AND MUNGED IN SOME
STANDARD WAY TO MANUFACTURE THE NEXT
INTEGER IN THE SERVENCE.

THE ENTIRE STATE OF THE RNG SHOULD NEVER

BE USED AS A "RANDOM" NUMBER. USE A FRACTION

OF IT, STARTING FROM THE LSB. MORE SIGNIFICANT

BITS WILL EVOLVE ACCORDING TO A MORE OBVIOUS PATTERN.

(4) WHERE TO START THE PSEUDO-RANDOM SERVENCE?

IT IS COMMONLY INITIALIZED FROM THE SYSTEM

ENTROPY CACHE (/dev/urandom), CPU clock, ETC.

* OR EQUIVALENTLY, VECTORS OF SMALLER INTEGERS

RNG, CONT.

- (5) A SEED (STATE VECTOR, OR PART THEREOF)
 MAY BE SAVED FOR FUTURE REFERENCE, OR
 VSED TO RE-INITIALIZE THE RNG.
 - THIS IS A HELPFUL WAY TO REPEAT

 EXACTLY THE SAME STRING OF PSEUDO-RANDOM

 NUMBERS AGAIN IN THE FUTURE. THIS

 REPEATABLE PANDOMNIESS IS OFTEN HELPFUL

 FOR CODE DEVELOPMENT, DEBUGGING, OR

 MAKING YOUR SCIENTIFIC RESULTS STRICTLY

 REPEATABLE.
 - BY SAVING THE SEED AFTER A STRING OF
 RAHDOM CALCULATIONS, YOU CAN PICK UP
 LATER WHERE YOU LEFT OFF, THIS IS
 HELPFUL ON SYSTEMS THAT AREN'T SO
 GOOD AT INITIALIZING THE RNG.
 - PERIODICITIES IN THE RNG ARE MOST

 EASILY SPOTTED BY FILLING A MANY-DIMENSIONAL

 CUBE WITH POINTS, AND LOOKING FOR CLUSTERING

 IN PLANES.

IN OTHER WORDS, RNG FLAWS BECOME APPARENT SOMETIMES IN HIGH-DIMENSIONAL MONTE CARLO WORK!

CAN WE DO BETTER THAN I/AN?

THERE ARE MANY APPROACHES TO SPEEDING CONVERGENCE FOR MONTE CIRCO METHODS.

SEE http://www-graphics.stanford.edu/papers/veach_thesis

- ONE CLASS OF METHODS REDUCES THE VARIANCE BY DISTRIBUTING THE POINTS SOMEWHAT MORE UNIFORMLY IN THE VOLUME.
 - · QUASI-RANDOM NUMBER SEQUENCES
 E.G. SOBOL NUMBERS
 - · STRATIFIED SAMPLING
 - Q: If MORE UNIFORMITY HELPS, WHY NOT USE A REGULAR GRID, AND FORGET THE RANDOM STUFF?
 - A: THAT PUTS US BACK AT THE TRAPEZOID OR MORSE!
 - THEN HOW MIGHT SUME DEGREE OF UNIFORMITY
 HELP, AND WHAT MIGHT WE MEAN BY THAT?

STRATIFIED SAMPLING

- OTHE IDEA IS TO DIVIDE THE DOMAIN INTO SUBAREAS, AND GET THE VOLUME OF EACH BY REGULAR M-C.
- · WORKS BEST IF THE INTEGRAND IS RELATIVELY CONSTANT WITHIN EACH SUBAREA.
- · WE CAN COME CLOSE TO THE IDEAL BY DIVIDING INTO N ERVAL SUBAREAS:

1D EXAMPLE:
$$n = 4$$
, 4 SUBAREAS, 1 POINT RANDOMLY PLACED IN EACH,

$$h = (b-a)/n$$

$$x = a + h(0:n-1) + h * rand(n)$$
STARTING [0, h, 2h, ... b-h] n UNIFORMLY DISTRIBUTED RANDOM DEVIATES BETWEEN n UNIFORMLY DISTRIBUTED n RANDOM DEVIATES BETWEEN n DISTRIBUTED n RANDOM DEVIATES BETWEEN n DISTRIBUTED n RANDOM DEVIATES n RANDOM DEVIATE

KEY POINTS

- PROBABLE.
- · POINTS DO NOT CLUSTER TOGETHER.
- · MOST SUBAREAS ARE VERY UNIFORM, BELAUSE WE'VE MADE THEM AS SMALL AS POSSIBLE.

HOW & WHY DOES STRATIFIED SAMPLING HELP?

CONSIDER THE V-SPHERE, WITH A 3D GRID OF MITHIN THEIR IN-VOLUME SUBDOMAINS.

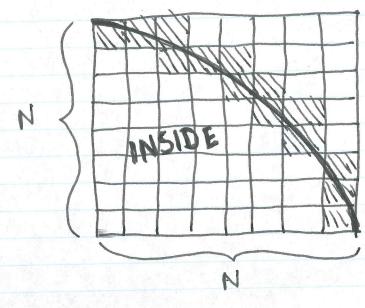
THERE ARE 3 KINDS OF SUBDOMAINS:

- THOSE FULLY INSIDE THE HYPERSPHERE, EACH CONTRIBUTES ho OF VOLUME, WHICH IS EXACTLY CORRECT. THERE ARE NEARLY PM. OF THESE.
- 2) THOSE FULLY OUTSIDE THE HYPERSPHERE, WHICH DON'T CONTRIBUTE ANYTHING.

ALL THE RANDOM ERROR COMES FROM THE LAST CATEGORY:

3) THOSE STRADDLING THE SURFACE OF THE HYPERSPHERE, SINCE THE SURFACE HAS V-1 DIMENSIONS, THE NUMBER OF STRADDLING SUBDOMBING SCALES AS $\mathcal{N}^{(v-1)/v}$. THE FRACTIONAL ERROR IN THE SURFACE CONTRIBUTION TO THE VOLUME IS $\frac{\Delta V_{\text{SURF}}}{V_{\text{SURF}}} = \frac{1}{\sqrt{n^{(v-1)/v}}} = \frac{n^{-(v-1)/2v}}{\sqrt{n^{(v-1)/v}}}$ THIS IS ONLY A FRACTION OF THE TOTAL VOLUME; $\frac{V_{\text{SURF}}}{V_{\text{TOT}}} = \frac{n^{-(v-1)/2v}}{\sqrt{n^{(v-1)/v}}} = \frac{n^{-(v-1)/2v}}{\sqrt{n^{(v-1)/v}}}$ SO THE RANDOM ERROR IN THE TOTAL VOLUME IS: $\frac{\Delta V_{\text{TOT}}}{\sqrt{n^{(v-1)/v}}} = \frac{\Delta V_{\text{SURF}}}{\sqrt{n^{(v-1)/v}}} = \frac{n^{-(v+1)/2v}}{\sqrt{n^{(v-1)/v}}}$

 $\frac{\Delta V_{ToT}}{V_{ToT}} = \frac{\Delta V_{SVRF}}{V_{TOT}} = \frac{\Delta V_{SVRF}}{V_{SVRF}} = \frac{V_{SVRF}}{V_{TOT}} = \frac{1}{N} \frac{-(\nu+1)/2\nu}{V_{TOT}}$ $\frac{\Delta V}{V} = \frac{1}{N} \frac{-(\frac{1}{2} + \frac{1}{2\nu})}{V_{TOT}} + \frac{1}{N} \frac{$



EDGE

$$\frac{3D}{N = N^2} = N^3$$

$$\frac{3D}{N = N^3}$$

$$\frac{2D}{N} = N^3$$

$$\frac{3D}{2N} = N^3$$

$$2N^2 = N^3$$

$$2N_{EDGE} \propto N^3$$

$$\frac{\nu - DIMENSIONS}{2}$$
 $2 = N^{\nu}$
 $2 = N^{\nu}$

:
$$n_{\text{EDGE}} \propto N^{\nu-1} = (n^{\frac{1}{\nu}})^{\nu-1} = n^{(\nu-1)/\nu}$$

FRACTIONAL ERROR



$$\frac{V_{SURF}}{V_{TOT}} = \frac{\mathcal{N}_{SURF}}{\mathcal{N}} \propto \frac{\mathcal{N}}{\mathcal{N}}$$